DATOMIC VS. CRUX
AND WHY IT MATTERS

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OUTUBRO 1982
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Micro Sistemas

A PRIMEIRA REVISTA BRASILEIRA DE MICROCOMPUTADORES

VisiCalc
Tecnologia aberta no DEL

SORTs comparados

Conheça o 6502
Impressoras

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12490

BYTE®

the small systems journal

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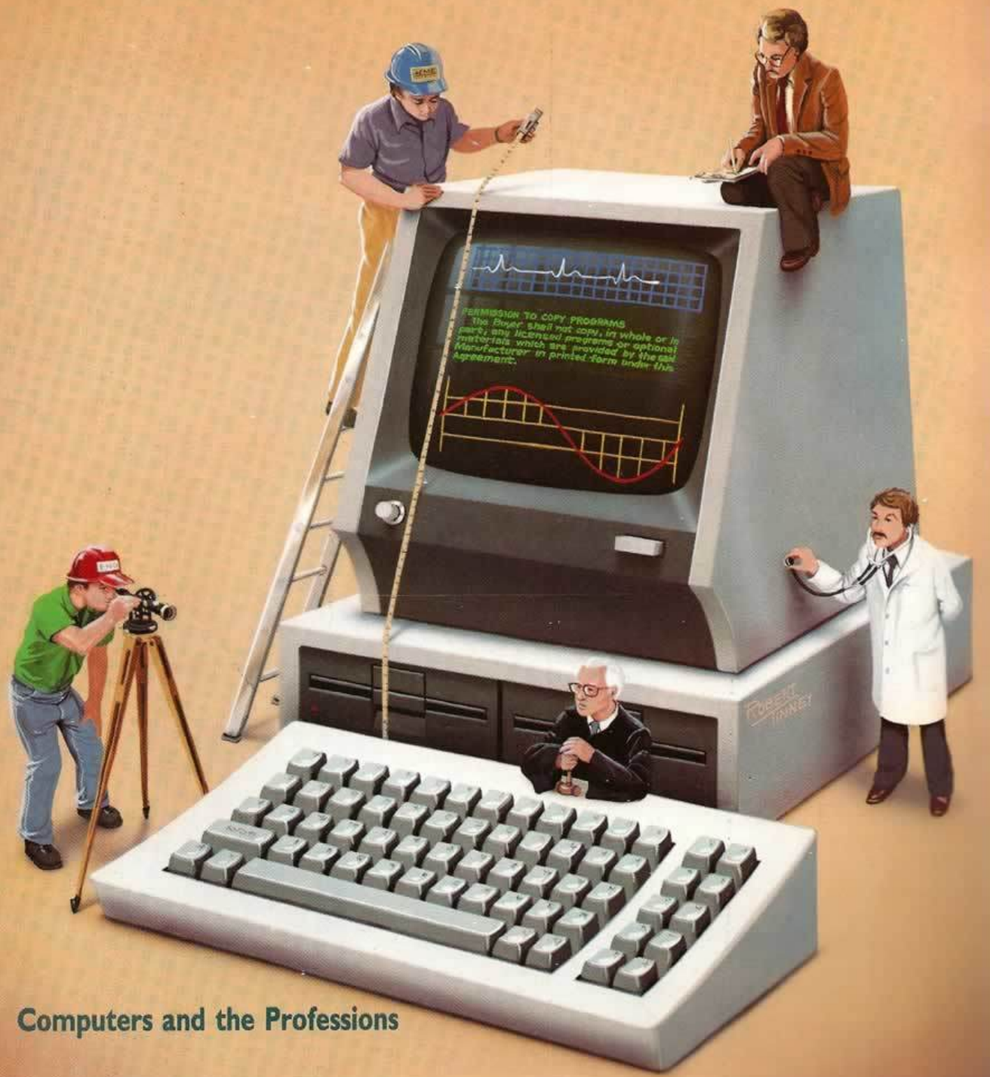

Computers and the Professions

```
670    INPUT "      ESTA TUDO CORRETO (S/N) ";ES
671    IF S$="N" THEN GOTO 620
672    IF SNOME$()"" THEN MAT$(I,1)=SNOME$
673    IF NOME$()"" THEN MAT$(I,2)=NOME$
674    IF ENDE$()"" THEN MAT$(I,3)=ENDE$
675    IF TEL$()"" THEN MAT$(I,4)=TEL$
676    MD%=1
677 NEXT I
678 IF MD%=1 GOTO 683
679 PRINT
680 PRINT "      SOBRENOME INEXISTENTE"
681 FOR K=1 TO 1000
682 NEXT K
683 RETURN
700 REM ------------------------------------------------( OPCAO PESQUISA )
701 PI=0
702 E%=0
703 HOME
704 PRINT "              PESQUISA"
705 VTAB 6
706 INPUT "    ENTRE COM O SOBRENOME ";SNOME$
707 SNOME$=SNOME$+SPACE$(N1%-LEN(SNOME$))
708 FOR I=2 TO NMAX%+1
709    IF SNOME$()MAT$(I,1) OR  MAT$(I,5)()"I"  THEN GOTO 732
710    E%=1
711    HOME
712    PRINT "            PESQUISA"
713    VTAB 7
714    PRINT
715    PRINT "SOBRENOME ";MAT$(I,1)
716    PRINT
717    PRINT "NOME      ";MAT$(I,2)
718    PRINT
719    PRINT "ENDERECO  ";MAT$(I,3)
720    PRINT
721    PRINT "TELEFONE  ";MAT$(I,4)
722    PRINT
723    VTAB 21
724    INPUT "    ESTE E' O REGISTRO PEDIDO (S/N) ";SN$
725    IF SN$="N" OR SN$="S" GOTO 729
726    PRINT
727    INPUT "    ENTRE S OU N ";SN$
728    GOTO 725
729    IF SN$="N" GOTO 732
730    PI=1
731    GOTO 733
732 NEXT I
733 IF E%=1 GOTO 739
734 PRINT
735 PRINT "     SOBRENOME INEXISTENTE"
736 FOR K=1 TO 1000
737 NEXT K
738 RETURN
739 IF PI=1 GOTO 745
740 IF PI()0 OR E%()1 GOTO 745
741 PRINT
742 PRINT "    NAO EXITE MAIS NENHUM ";SNOME$
743 FOR K=1 TO 1000
744 NEXT K
745 RETURN
800 REM ------------------------------------------------( OPCAO LISTAGEM )
801 REM
802 REM ------------------------------------( TELA DE OPCOES PARA CLASSIFICACAO )
803 REM
804 HOME
805 PRINT "              LISTAGEM"
806 VTAB 9
807 PRINT "    ESCOLHA A CHAVE DE CLASSIFICACAO"
808 PRINT
809 PRINT "           SOBRENOME (1)"
810 PRINT "           NOME      (2)"
811 PRINT "           ENDERECO  (3)"
812 PRINT "           TELEFONE  (4)"
813 PRINT
814 INPUT "    ENTRE COM A OPCAO ";OP%
815 IF OP%>=1 AND OP%<=4 GOTO 821
816 PRINT
817 PRINT "     CHAVE INEXISTENTE"
818 FOR K=1 TO 1000
819 NEXT K
820 GOTO 804
821 IF OP%()1 GOTO 824
822 NCHAR%=N1%
823 INIC%=1
824 IF OP%()2 GOTO 827
825 NCHAR%=N2%
826 INIC%=11
827 IF OP%()3 GOTO 830
828 NCHAR%=N3%
829 INIC%=41
830 IF OP%()4 GOTO 833
831 NCHAR%=N4%
832 INIC%=81
833 J=1
834 FOR I=2 TO NMAX%+1
835    IF MAT$(I,5)()"I" GOTO 843
836    A$=MAT$(I,1)+SPACE$(N1%-LEN(MAT$(I,1)))
837    B$=MAT$(I,2)+SPACE$(N2%-LEN(MAT$(I,2)))
838    C$=MAT$(I,3)+SPACE$(N3%-LEN(MAT$(I,3)))
839    D$=MAT$(I,4)+SPACE$(N4%-LEN(MAT$(I,4)))
840    E$=MAT$(I,5)
841    R$(J)=A$+B$+C$+D$+E$
842    J=J+1
843 NEXT I
844 NREG%=J-1
845 REM ------------------------------------------------( OPCAO CLASSIFICACAO )
846 GOSUB 900
847 FOR I=2 TO NTOTAL%+1
848    MAT$(I,1)=MID$(R$(I-1),1,N1%)
849    MAT$(I,2)=MID$(R$(I-1),N1%+1,N2%)
850    MAT$(I,3)=MID$(R$(I-1),N1%+N2%+1,N3%)
851    MAT$(I,4)=MID$(R$(I-1),N1%+N2%+N3%+1,N4%)
852    MAT$(I,5)=MID$(R$(I-1),N1%+N2%+N3%+N4%+1,1)
853 NEXT I
854 REM ------------------------------------------------( LISTAGEM )
855 FOR I=2 TO NTOTAL%+1
856    HOME
857    PRINT "              LISTAGEM"
858    VTAB 9
859    PRINT "SOBRENOME  ";MAT$(I,1)
860    PRINT
861    PRINT "NOME       ";MAT$(I,2)
862    PRINT
863    PRINT "ENDERECO   ";MAT$(I,3)
864    PRINT
865    PRINT "TELEFONE   ";MAT$(I,4)
866    VTAB 23
867    INPUT "ENTRE RETURN PARA CONTINUAR ";AA$
868 NEXT I
869 RETURN
900 REM ------------------------------------------( SUBROTINA DE CLASSIFICACAO )
901 DEFINT A-J,L-Q
902 DEFINT S-Z
903 M=9
904 K$(0)=STRING$(NCHAR,0)
905 K$(NREG+1)=STRING$(NCHAR,127)
906 REM
907 FOR I=1 TO NREG
908    K$(I)=MID$(R$(I),INIC,NCHAR)
909 NEXT I
910 IF NREG<=M THEN GOTO 961
911 IP=1
912 P(IP,1)=0
913 P(IP,2)=0
914 L=1
915 S=NREG
916    I=L
917    J=S+1
918    KEY$=K$(L)
919    I=I+1
920 IF K$(I)<KEY$ THEN GOTO 919
921 J=J-1
```

CRUX & DATOMIC
REASONS WHY YOU SHOULD CARE

REASONS WHY YOU SHOULD CARE

1. Immutable Database
2. Query Like a Ninja
3. Unbundle Your Database

REASONS WHY YOU SHOULD CARE

1. Immutable Database
2. Query Like a Ninja
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IMMUTABILITY EVERYWHERE

» Haskell, Clojure, Erlang...

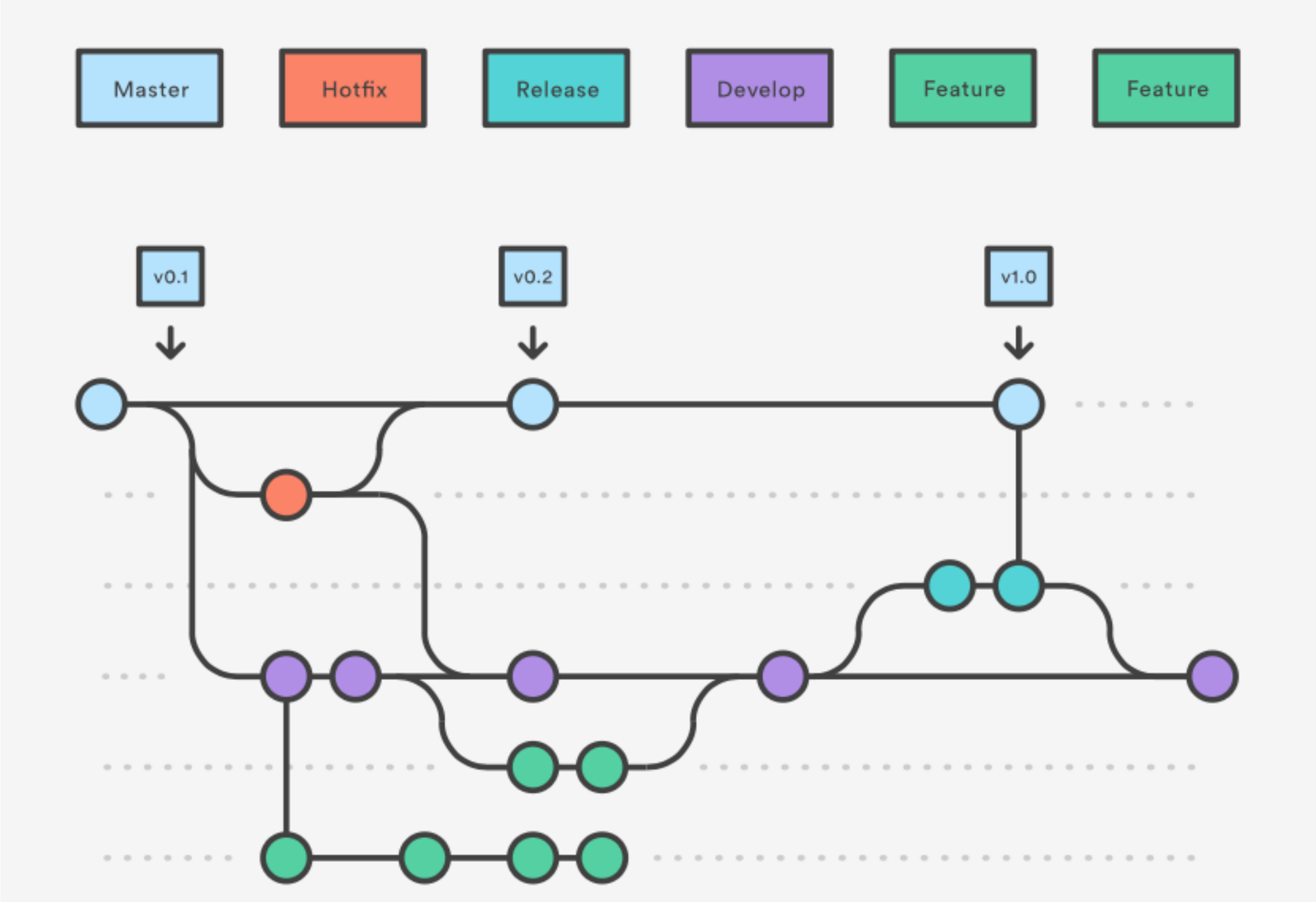
» Immutable.js

» CloudFormation, Terraform

» Nix OS

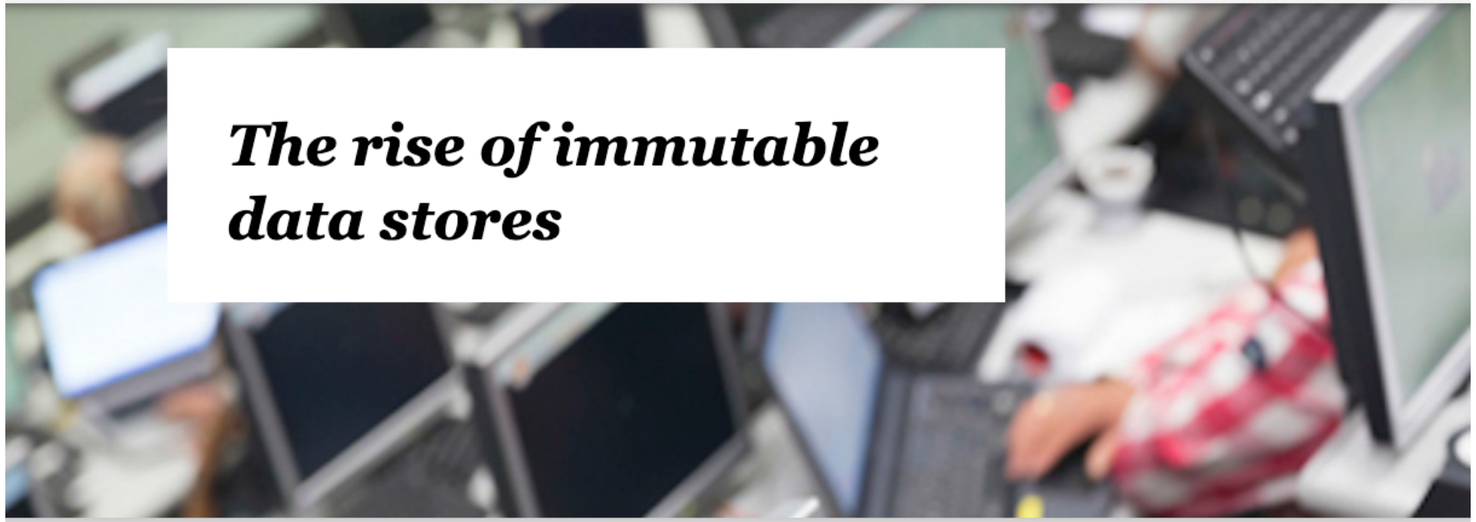
MYTH:

DATA IS EPHEMERAL


Master
Hotfix
Release
Develop
Feature
Feature
v0.1
v0.2
v1.0

ACCRETING FACTS OVER TIME

pwc | Next in Tech



The rise of immutable data stores

by Alan Morrison

September 8, 2015

Tags: Data NoSQL Databases Unstructured Data

Some innovators are abandoning long-held database principles. Why?

The website for Room Key, a joint venture of six hotel chains to help travelers find and book lodging, collects data from as many as 17 million pages per month, records an average of 1.5 million events every 24 hours, and handles peak loads of 30 events per second. To process that onslaught of complex information, its database records each event without waiting for some other part of the system to do something first.

REASONS WHY YOU SHOULD CARE

1. Immutable Database
2. Query Like a Ninja
3. Unbundle Your Database

CRITIQUE OF SQL

» lack of proper orthogonality

» lack of compactness

» lack of consistency

» poor system cohesion

MAY 09, 2019

We Can Do Better Than SQL

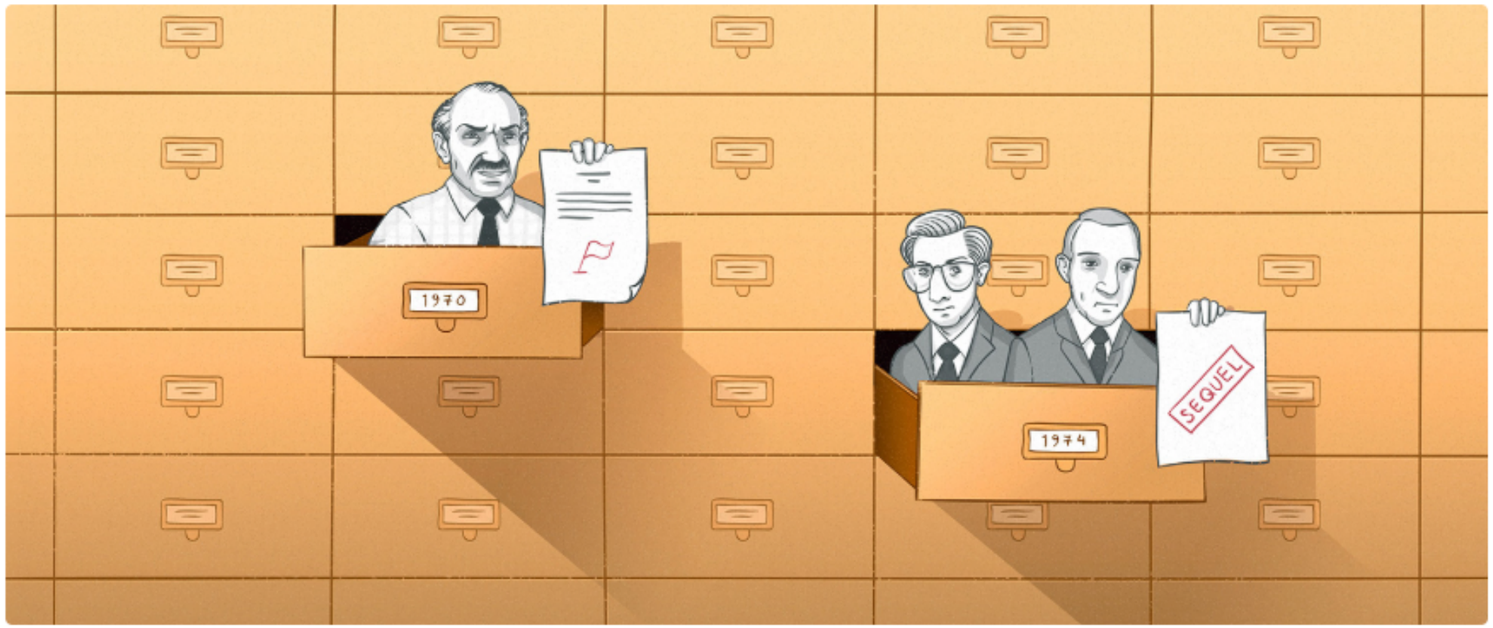
Elvis Pranskevichus

all blog posts

A Brief History of the Relational Model and SQL
Critique of SQL
Why Does All This Matter?
EdgeQL: Query Language Evolution
Final Words

tweet

subscribe



The questions we often hear are "Why create a new query language?" and "What's wrong with SQL?". This post contains answers to both.

Before we begin, let's overview some of the history of how the relational model came to be, and how SQL was created.

DATALOG (BOSS SLIDE)

Datalog is a declarative query language with roots in logic programming that combines facts and rules to achieve the same power as relational algebra recursion.

```
SELECT greatest_songs.name
FROM artists
INNER JOIN greatest_songsa
ON artists.id = greatest_songs.artist_id
WHERE artists.name = "Luan Santana";
```

```
[:find ?s
 :where
 [?s :greatest-song/artist ?p]
 [?p :artist/name "Luan Santana"]]
```

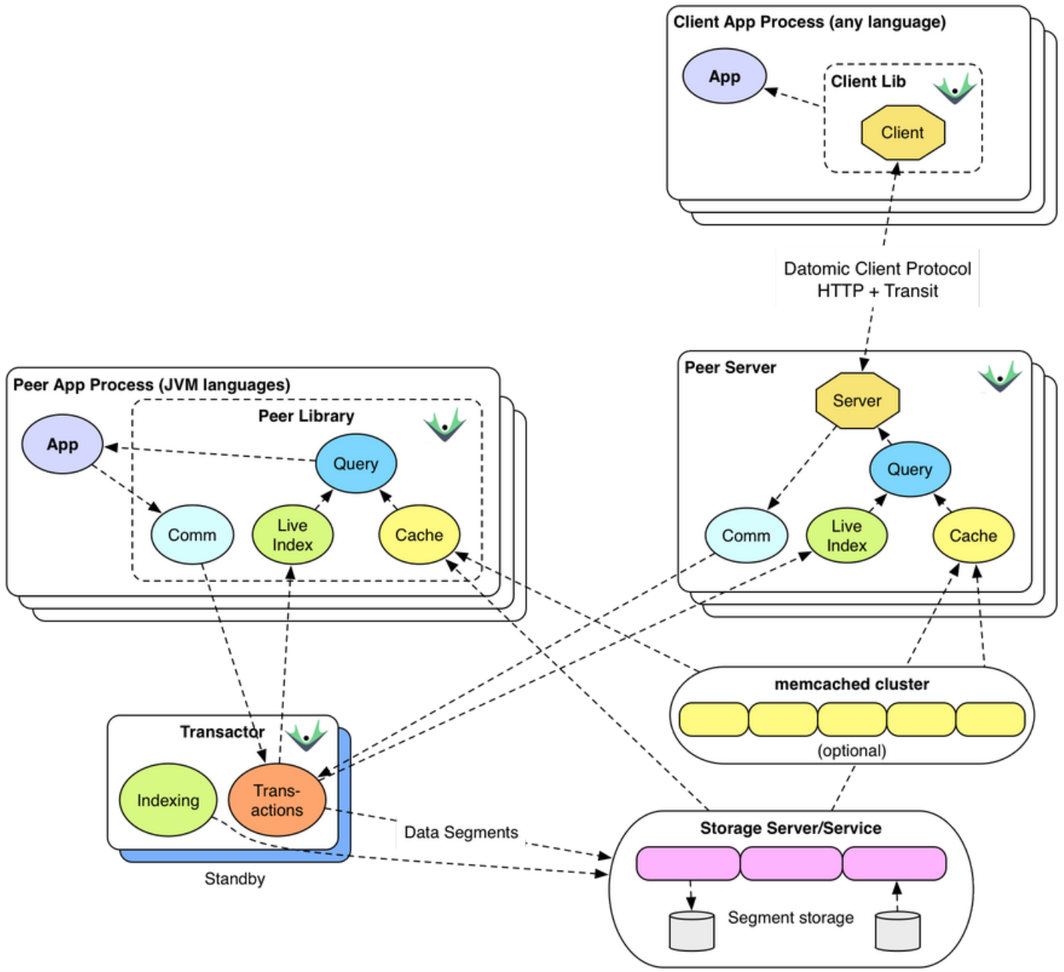
HTTP://WWW.LEARNDATALOGTODAY.ORG

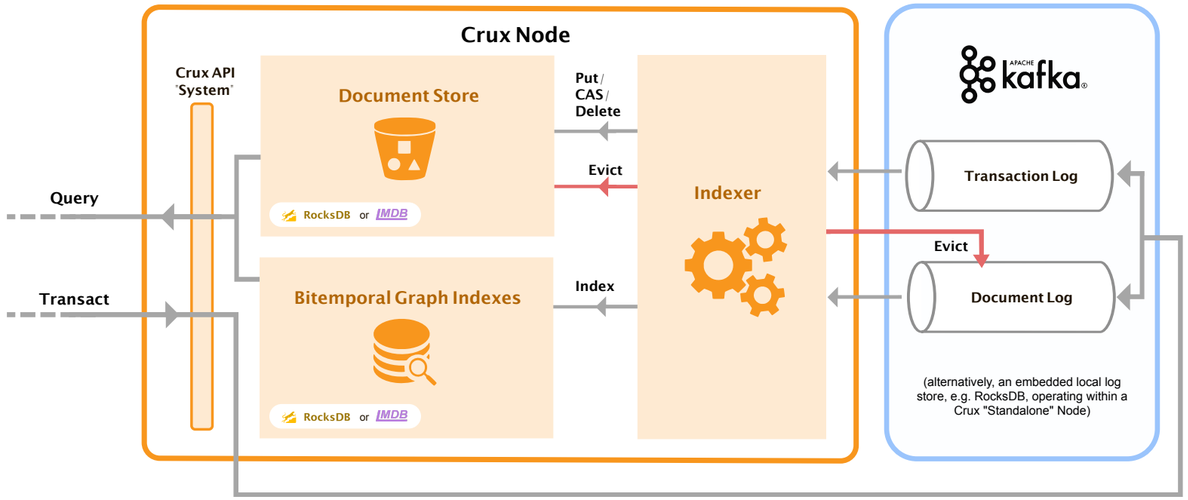
REASONS WHY YOU SHOULD CARE

1. "Immutable" Database
2. Query Like a Zen Master
3. The Unbundled Database

THE UNBUNDLED DATABASE (BOSS SLIDE)

“What do we have to gain from turning the database inside out? Simpler code, better scalability, better robustness, lower latency, and more flexibility for doing interesting things with data.”
Martin Kleppmann


Client App Process (any language)
App
Client Lib
Client
Datomic Client Protocol
HTTP + Transit
Peer App Process (JVM languages)
Peer Library
App
Query
Comm
Live Index
Cache
Peer Server
Server
Query
Comm
Live Index
Cache
memcached cluster
(optional)
Transactor
Indexing
Trans-actions
Standby
Data Segments
Storage Server/Service
Segment storage


Crux Node
Crux API 'System'
Query
Transact
Document Store
RocksDB or IMDB
Bitemporal Graph Indexes
RocksDB or IMDB
Put / CAS / Delete
Evict
Index
Indexer
kafka
Transaction Log
Evict
Document Log
(alternatively, an embedded local log store, e.g. RocksDB, operating within a Crux "Standalone" Node)

"TURNING THE DATABASE INSIDE OUT"

"Turning the database inside out with Apache Samza" by Martin Kleppmann

115,453 views 1.2K 7 SHARE SAVE

https://www.youtube.com/watch?v=fU9hR3kiOK0

REASONS WHY YOU SHOULD CARE

1. Immutable Database
2. Query Like a Ninka
3. Unbundle Your Database

C R U X

Datomic

THE 10 CRITICAL THINGS WHEN TALKING ABOUT DATOMIC AND CRUX

1. LICENSING MODEL

1. LICENSING MODEL

» Crux: Open Source (MIT)

» Datomic: Proprietary Only

2. TEMPORAL MODEL

2. TEMPORAL MODEL

» Crux: Bitemporal

» Datomic: Unitemporal

DATOMIC'S TEMPORAL MODEL

- "A datom is an immmutable, point-in-time fact!"
  - Entity
  - Attribute
  - Value
  - Time
- EAVT

Entity	Attribute	Value	Time
`123`	`:person/likes`	`"pizza"`	`1978-07-08T20:19:03.176-00:00`

Entity	Attribute	Value	Time
`123`	`:person/likes`	`"pizza"`	`1978-07-08T20:19:03.176-00:00`
`123`	`:person/name`	`"John"`	`1978-07-08T20:19:03.176-00:00`
`123`	`:contact/phone`	`"+1 (305) 555-5524"`	`2009-03-01T14:24:32.074-00:00`
`123`	`:contact/phone`	`"+1 (646) 341-3367"`	`2017-10-25T08:22:06.165-00:00`

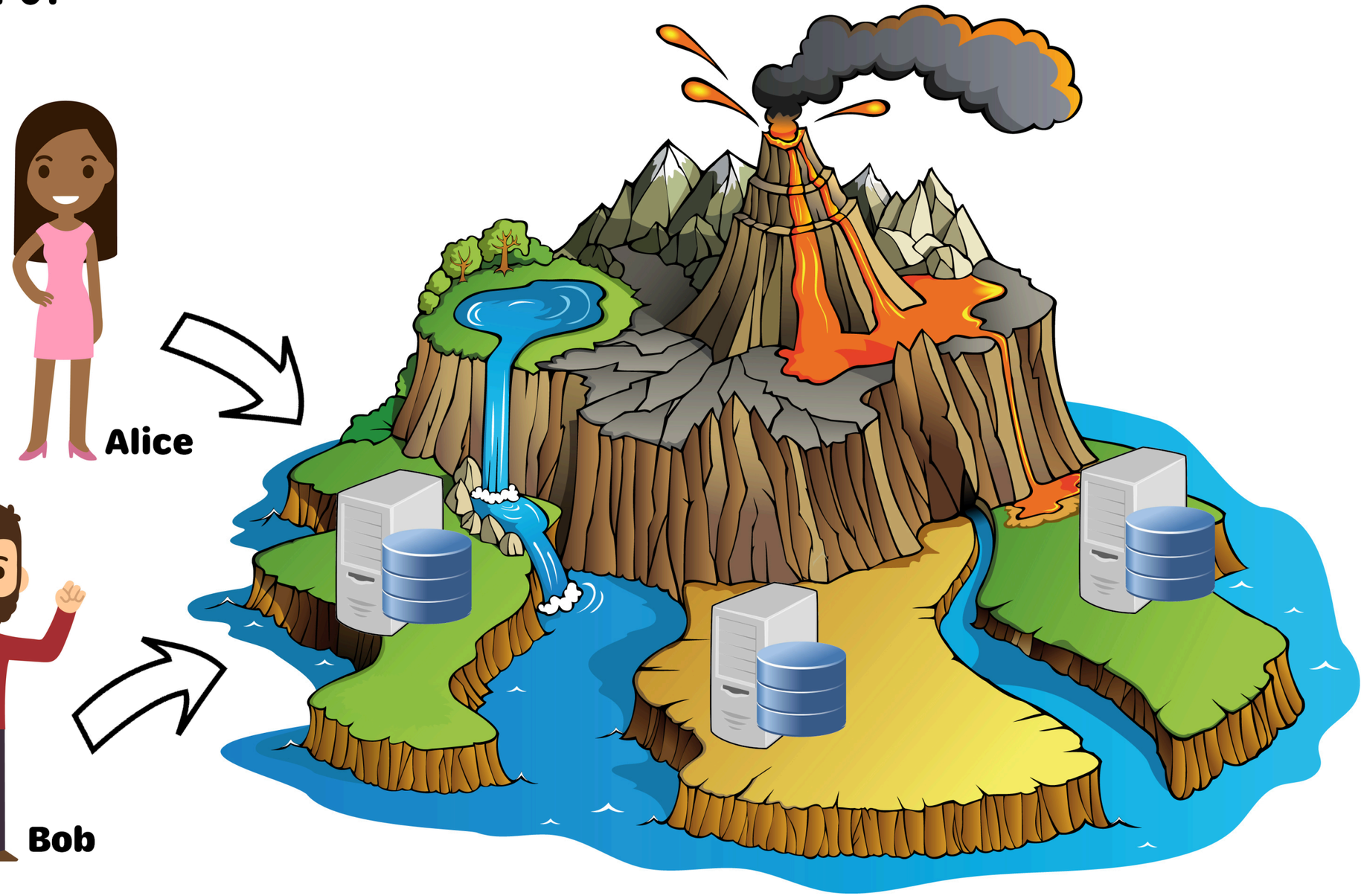
CRUX'S BITEMPORAL MODEL

» Transaction Time

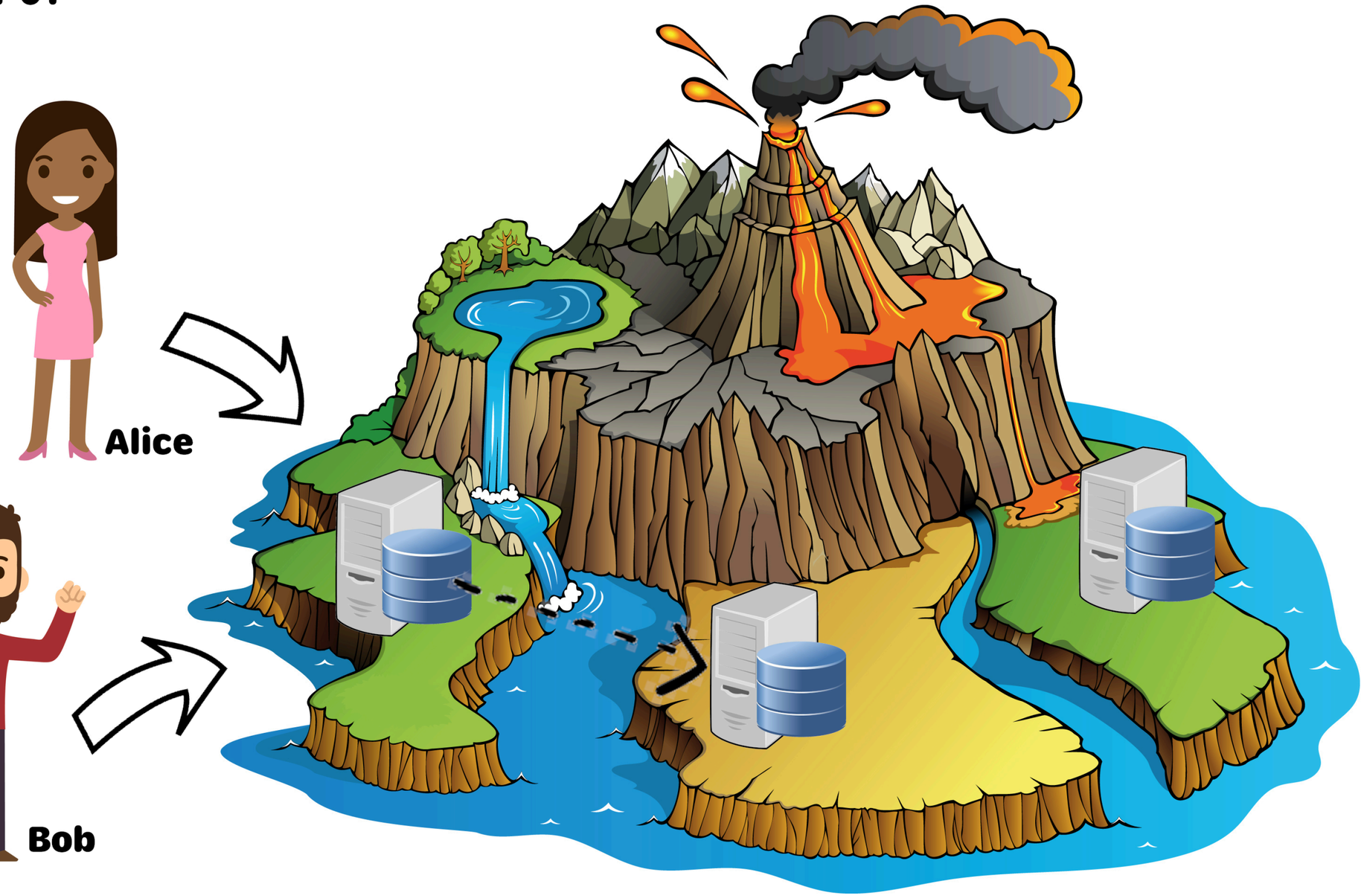
» Valid Time

2019-01-01

2019-01-01

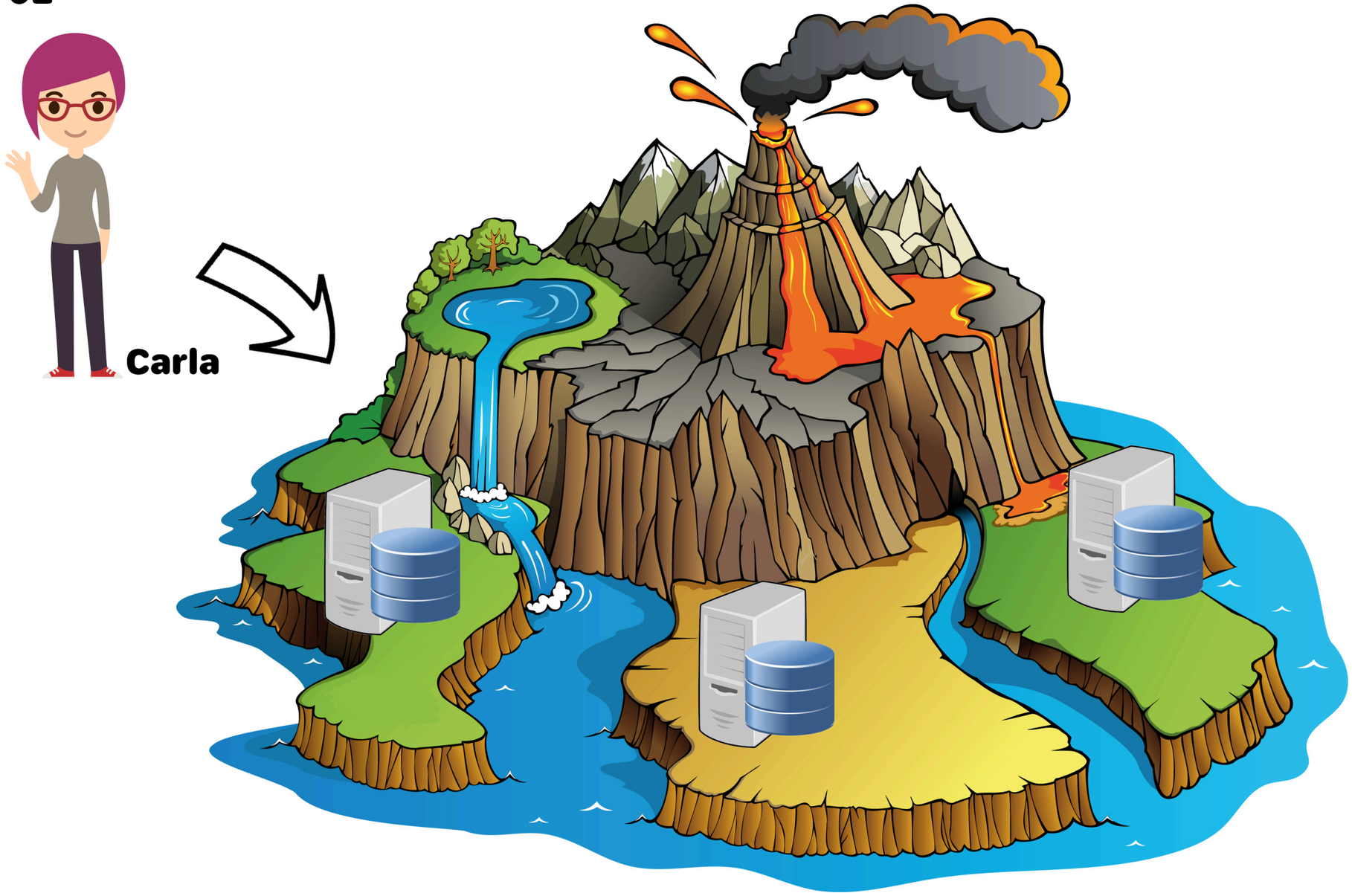


2019-01-01

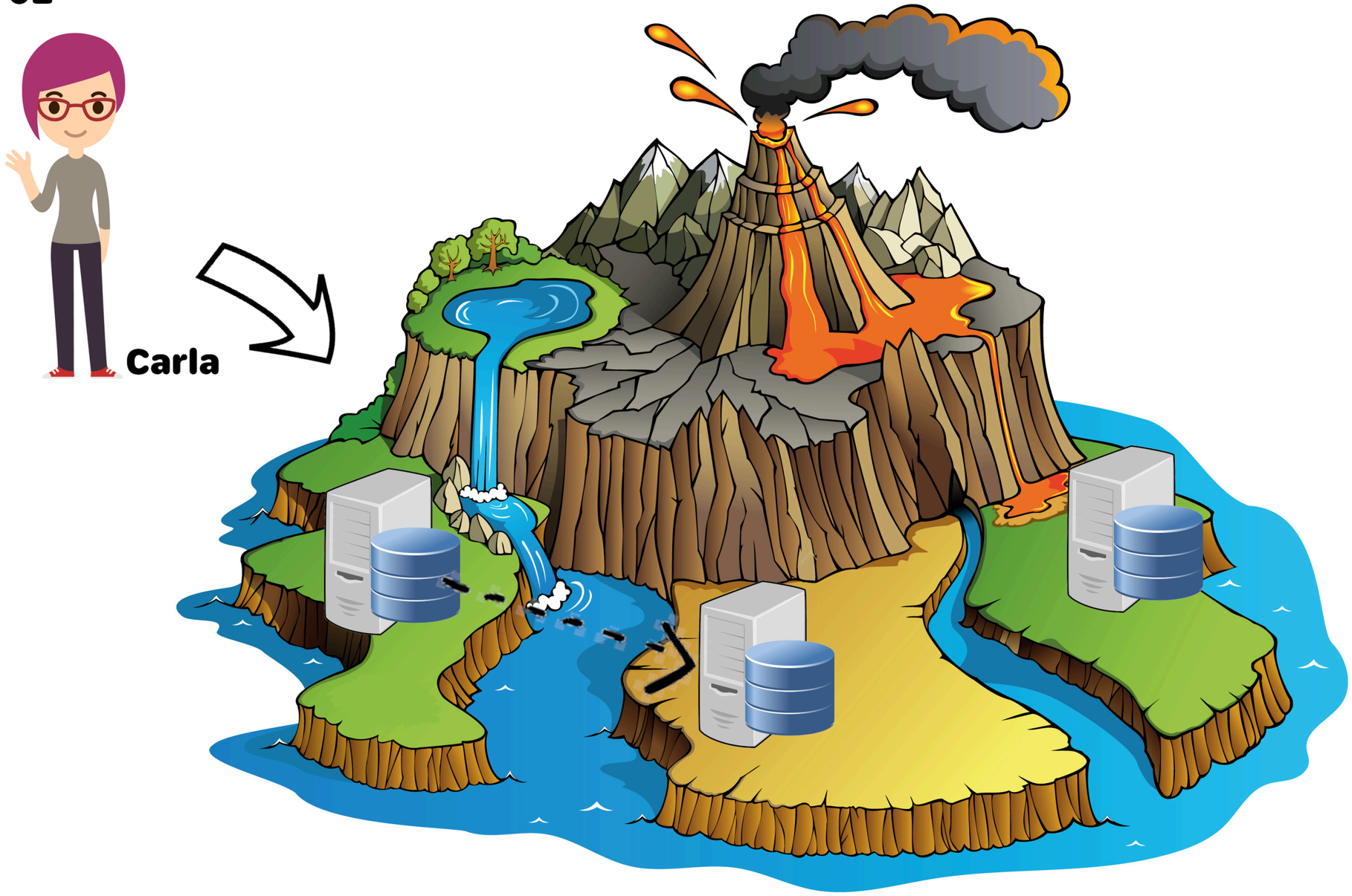


2019-01-02

2019-01-02



2019-01-02



2019-01-03

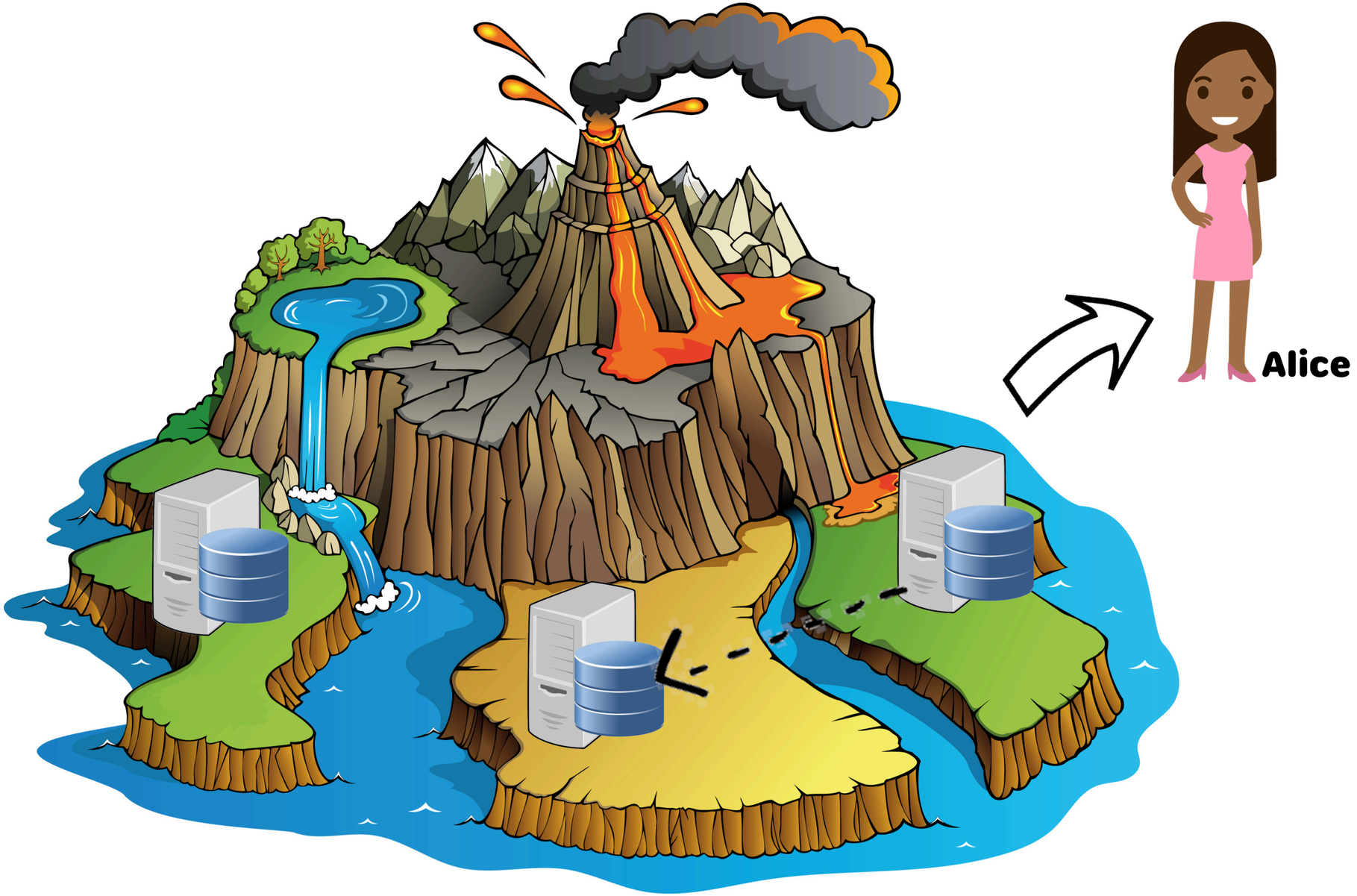
2019-01-03

2019-01-03

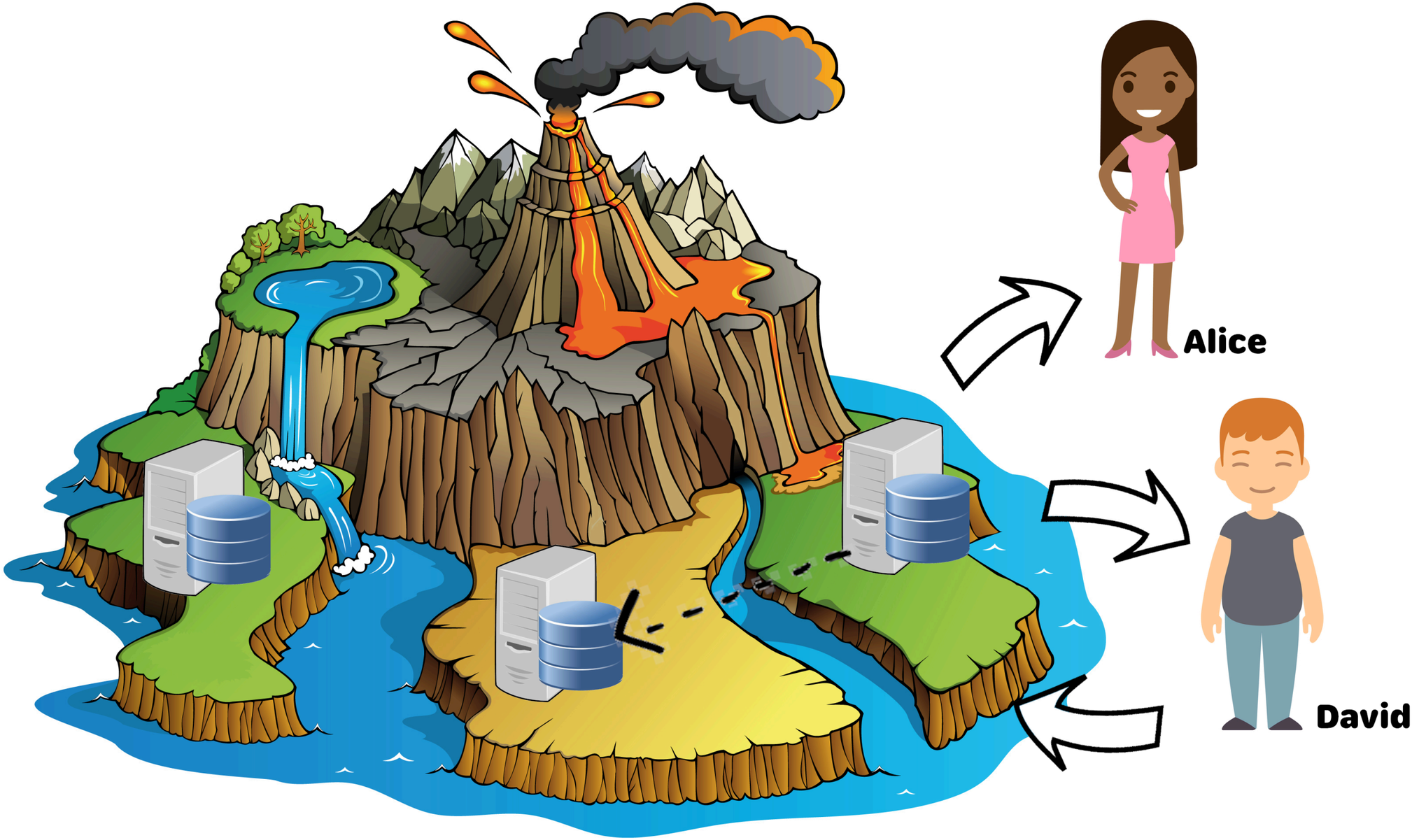
2019-01-04

2019-01-04

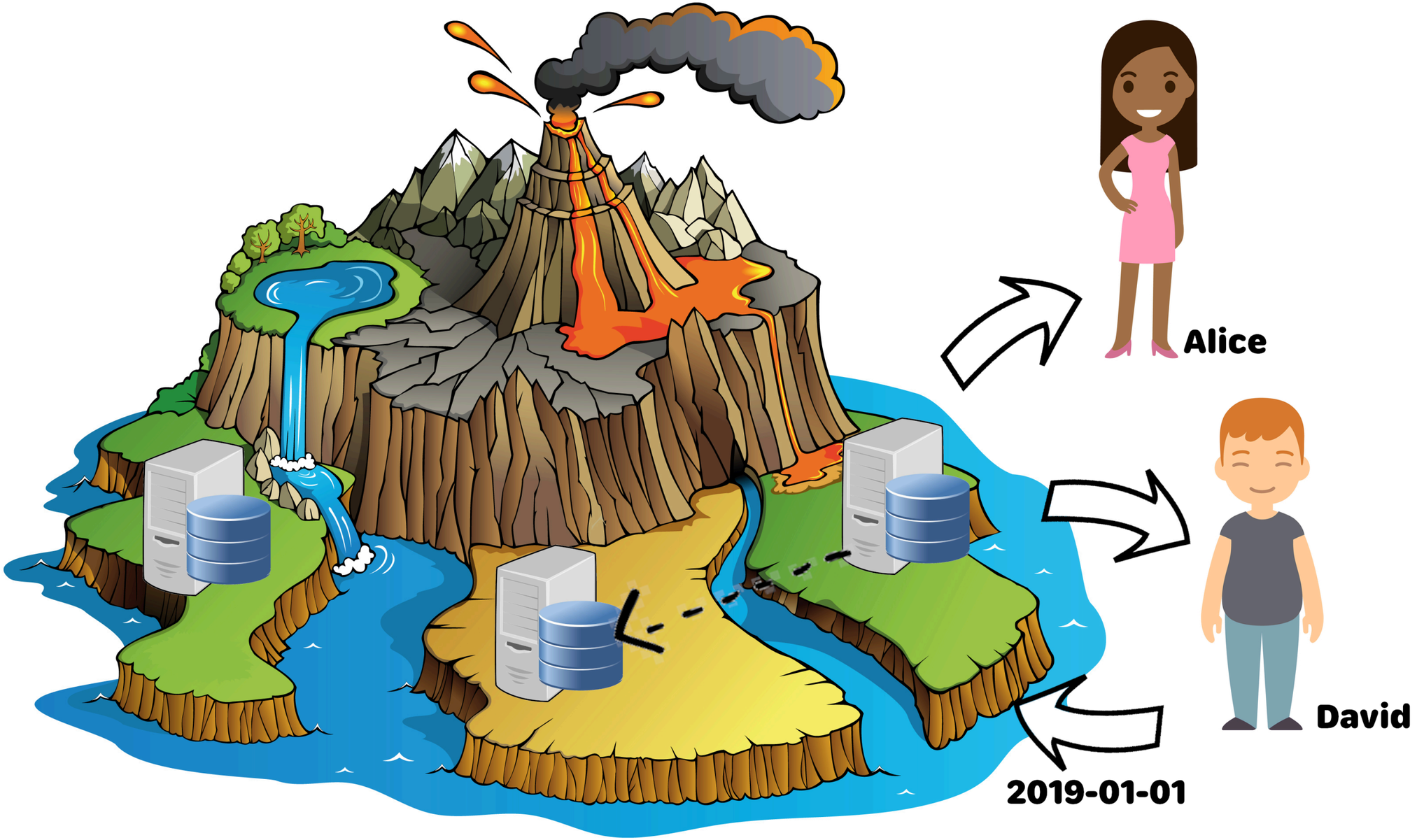
2019-01-04



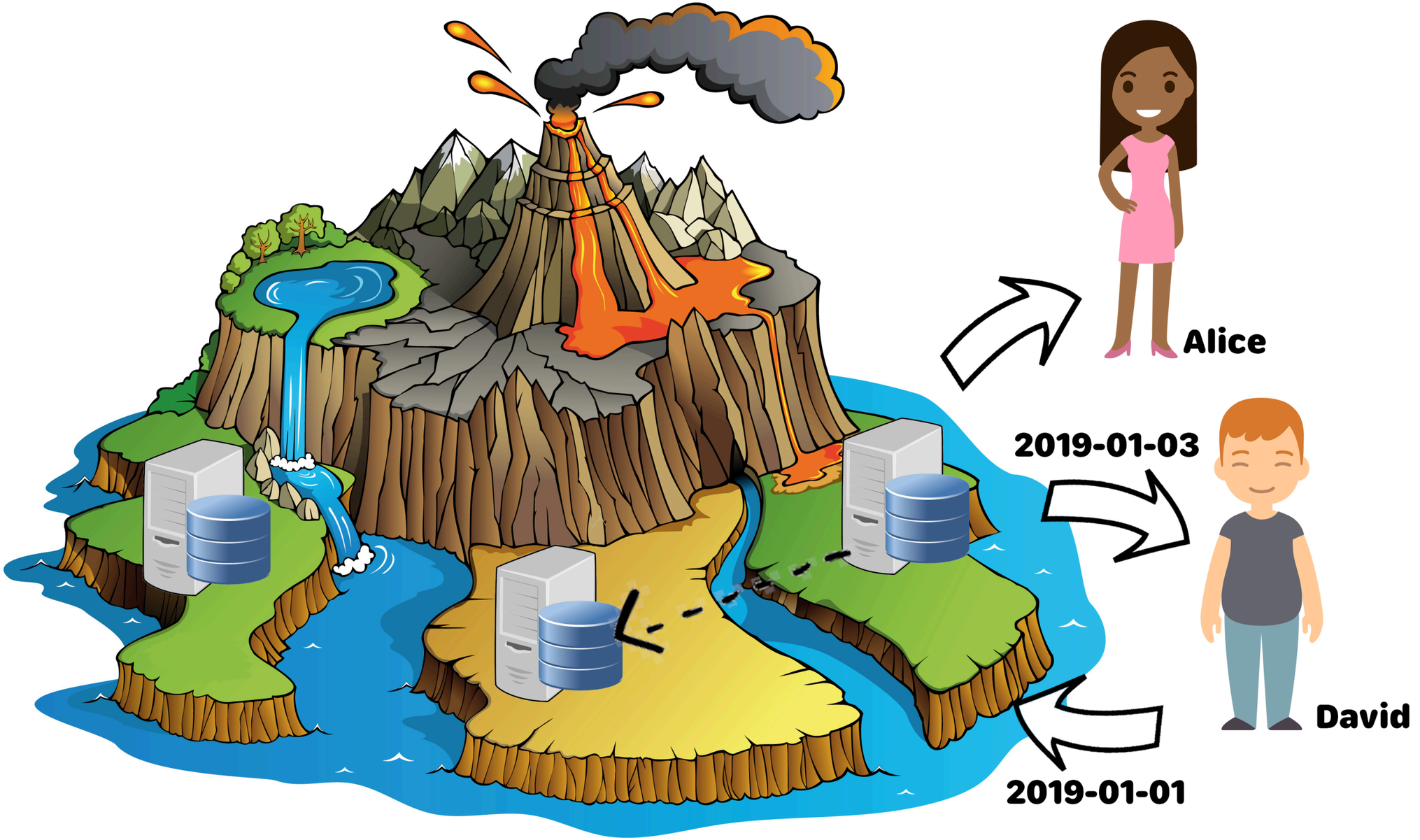
2019-01-04



2019-01-04



2019-01-04



WHO'S ON THE ISLAND ON THE 2ND?

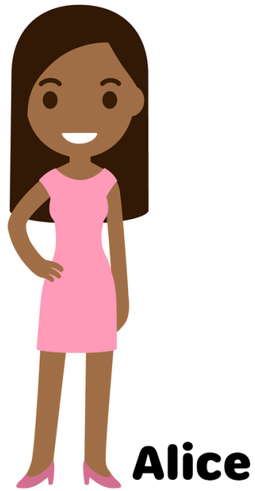
WHO'S ON THE ISLAND?

» "as of" - Transaction time

» "as at" - Valid time

WHO'S ON THE ISLAND ON THE 2ND? (AS OF THE 3RD)

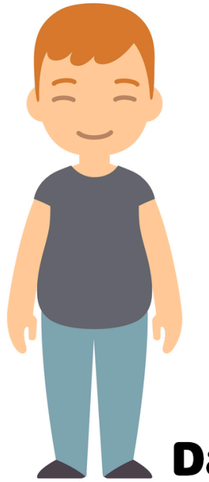
WHO'S ON THE ISLAND ON THE 2ND? (AS OF THE 4TH)



Alice

Bob

Carla



David

BITEMPORALITY (BOSS SLIDE)

» Ingest out-of-order temporal data from upstream timestamping systems

» Reconcile temporal data across eventually consistent systems

» Dealing with backlog corrections

AUDIT ALL DATA CHANGES AND PERFORM DATA FORENSICS

3. STORAGE MODEL

3. STORAGE MODEL

» Crux: Document-based

» Datomic: Entity-Attribute-Value-based

CRUX BEING DOCUMENT-BASED MEANS

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1. Transactions: apply to the whole doc

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2. Whole document gets changed on transactions

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3. Only root-level nodes are indexed

CRUX BEING DOCUMENT-BASED MEANS

1. Transactions: apply to the whole doc
2. Whole document gets changed on transactions
3. Only root-level nodes are indexed
4. Document relationships can be tricky

4. SCHEMA MODEL

4. SCHEMA MODEL

» Crux: No Schema Definition/Enforcement

» Datomic: Required Schema Definition/Enforcement

CRUX HAVING NO SCHEMA MEANS

- No validation (anything goes)
  - Decoupled concerns
- Silver lining
  - No explicit schema evolution process needed
  - Good practice: implement your own validations

5. DEPLOYMENT MODEL

5. DEPLOYMENT MODEL

- Crux: Self Hosted
  - Confluent Cloud (KaaS)
  - Managed hosting available (via Juxt)
- Datomic:
  - Datomic On-prem: Self Hosted
  - Datomic Cloud: Self Hosted (AWS Cloud-based)

SOME IMPLICATIONS

- Both can be easily developed locally
  - run locally, self-host somewhere, or containerize it
- Crux locally can be a standalone (crux-core)
- Datomic Cloud requires an AWS connection for development

6. PRIVACY MODEL

What is the right to erasure?

Under Article 17 of the GDPR individuals have the right to have personal data erased. This is also known as the ‘right to be forgotten’. The right is not absolute and only applies in certain circumstances.

6. PRIVACY MODEL

» Crux: "Eviction" operation

» Datomic:

  » Datomic On-prem: "Excision" operation

  » Datomic Cloud: nothing available

7. QUERY DIFFERENCES

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1. Queries are not portable across

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2. Crux does not support Pull Query

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5. Datomic requires clause order

7. QUERY DIFFERENCES

1. Queries are not portable across
2. Crux does not support Pull Query
3. Datomic does not have Lazy Queries
4. Crux does not have sophisticated look up refs
5. Datomic requires clause order
6. Several other minor differences

8. TRANSACTION DIFFERENCES

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1. Transactions are not portable across

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8. TRANSACTION DIFFERENCES

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2. Crux does not support Components
3. Crux does not support Transaction meta model

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5. Datomic's injection speed is bound to transactor

8. TRANSACTION DIFFERENCES

1. Transactions are not portable across
2. Crux does not support Components
3. Crux does not support Transaction meta model
4. Crux does not support Transaction functions
5. Datomic's injection speed is bound to transactor
6. And remember: doc-based vs. attribute-based

9. OTHER ARCHITECTURAL DIFFERENCES

9. OTHER ARCHITECTURAL DIFFERENCES

	CRUX	DATOMIC
Topology	Each Node Storage/ Index	Global Shared Storage/Index
Index Sharding	No Auto Sharding	Peers auto-shard based on working sets
Transactor model	Single Unpartitioned Kafka Topic	Explicit Transactor

10. OTHER "FREEBIES"

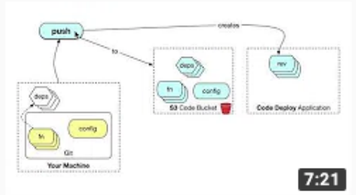
10. OTHER "FREEBIES"

- Datomic
  - Datomic Cloud
  - Datomic Ions
  - Datomic Cast

Rich Hickey on Datomic Ions, September 12, 2018

Clojure/nyc • 9.1K views • 11 months ago

The incidental complexity dragon never sleeps! In building **Datomic** Cloud we took on simplifying **Datomic** deployment, security ...



Datomic Ions in Seven Minutes

ClojureTV • 5.5K views • 1 year ago

Stuart Halloway introduces **Ions** for **Datomic** Cloud on AWS.

Declarative Domain Modeling for Datomic Ion/Cloud - Tiago Luchini

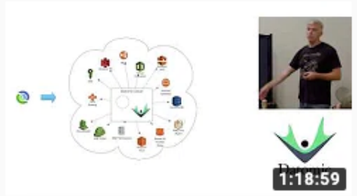
ClojureTV • 2.6K views • 8 months ago

What if we could build on top of **Datomic Ions** already easy-to-use, easy-to-deploy, setup and make some common scenarios even ...

Datomic Ions Hello World in 25 minutes

Tiago Luchini • 1.4K views • 1 year ago

I've set myself the target to see how long it would take to code and deploy a Hello World in **Datomic Ions**. We cover: - setting up a ...



Datomic Ions - Stuart Halloway

ClojureTV • 5.8K views • 1 year ago

Datomic Ions (https://goo.gl/XcEQNh) let you develop applications for the cloud by deploying your code to a running Datomic ...

Serverless-ish: Zero to App with Datomic Cloud and GraphQL - Chris Johnson Bidler

ClojureTV • 1.5K views • 8 months ago

The #serverless architecture pattern is taking the world by storm, and for good reasons. The principles of server less design allow ...

10. OTHER "FREEBIES"

- Crux
  - Open Source
  - I.e. Active Object Backend
  - Coming up: Crux console

CLOJURE IS GREATLY SERVED ON THE DATABASE SIDE

WITH EITHER CRUX OR DATOMIC:

1. Immutable Database
2. Query Like a Ninja
3. Unbundle Your Database

THANK YOU!

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